Research on the *k*-coverage local wireless network and its communication coordination mechanism design

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Abstract. Based on the analysis of the present situation and features of local wireless network, combined the networking technologies, we constructing the wireless network layout algorithm and sensor nodes model. Through delaminating iterative optimization algorithm, optimal layout of sensor data collection and satisfy recognition, quasi-geoids orientation, signal coverage optimization have been realized. Compared the triangular, square and hexagon scheme, and get a general scheme of multiple complete coverage by calculate the length of sides on the polygon. Monte Carlo algorithm is used in choosing the optimum proposal comparing among the regular triangle, square and the hexagon in different conditions. For locating network, analyze the advantages and disadvantages of triangle, square and hexagon schemes in deferent positioning accuracy of double coverage, we set different layout schemes to different situations. Application manual is provided for the constructor to refer to these table lookup to attain the right numerical value.

Keywords: *k*-coverage, local wireless network, communication coordination, Monte Carlo

1   Introduction

In the subject of internet of things, different kinds of information are transferred by sensor nodes. In order to improve the efficiency and stability of the transmission through networks, the layout of the nodes must be optimized. (Zhang Honghai and Hou Jennifer,2005; Yu Hongyi,2008)

Zigbee is one of the hardware equipment used in the wireless communication technology. It has many advantages, for example, the connection distance is shorter and the complexity is smaller, the power consumption is lower and so on. In the automatic control and remote control field, it is widespread. Many communicators are inset with Zigbee. It consists of three models: the response model; the data manipulation model; the communication model. The response model is to receive the information. The information is transferred through the communication model. With such information, we can attain the location of the moving nodes. (Duc A. Tran，Usman Khan，2006)

2   Methods and Algorithms

2.1   The network evaluation function

When to evaluate a network, we often consider the following two factors: the network performance and the cost. The network performance includes the stability in transfer information and the accuracy in location. It depends on the overlapping area. The cost depends on the number of nodes. The more nodes there are the higher cost and power consumption.

When the target area is k-coverage, the size of area covered is:

$$S_{be\,covered} = ab \qquad (1)$$

If the number of nodes in this area is n, the coverage area is:

$$S_{all} = n\pi R^2 \qquad (2)$$

Where: a is the length of the target area,

b is the width of the target area

R is the effective communication radius.

At the same time, every point in the target area has the very same information collection accuracy.

Definition 1: the coverage rate of sensor η

This passage comes up with the concept of coverage rate. It means the ratio between the size of target area and the sum of all the coverage area provided by the nodes.

$$\eta = \frac{\bigcup\limits_{i=1}^{n} S_i}{\sum\limits_{i=1}^{n} S_i} \tag{3}$$

It can reflect on the utilization of the coverage area. The larger η, the better the net performance is. At the same time, it can reflect the number of nodes. The fewer the nodes are, the lower the cost is. In another words, the lower the coverage rate, the better the layout of the network.

In a 1-coverage network, the calculation function is defined:

$$g_1 = (S_{cover} - S_{becovered}) / S_{becovered} = \frac{1}{\eta} - 1 \tag{4}$$

Through the discussion, the coverage rate is solved. In the 1-coverage network, it is 82.7%. In fact, it can't reach the perfect value 1.

In a k-coverage network, the covered area can be considered as:

$$kS_{becovered} = kab \tag{5}$$

All in all, in a k-coverage network, the calculation function is defined:

$$g_k = (S_{cover} - S_{becovered}) / S_{becovered} = \frac{k}{\eta} - 1 \tag{6}$$

2.2   The layout of nodes in 1-coverage network

To design the nodes layout, for a target area, the main problem is how to make the area is covered without blind spot, which means that the point can`t be covered. According to the simplification of the covered area—the circle, we can turn the problem into geometric. For a ab rectangle region, how to be fitted together can use the least circles and how many circles are needed. As no matter small the radius of the circle is, it impossible to cover the area without blind point, so we take the second best solution that the circle is replaced with its inscribed polygon and the node is located on the center of the polygon.

If there are x polygons put on the same vertex, and every n gon's interior angle

is:

$$(n-2)*180° \, / \, n \qquad (7)$$

Then:

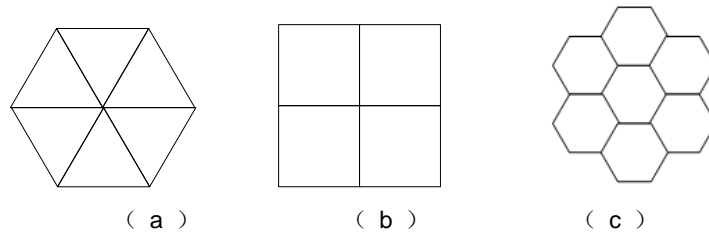$$x \frac{(n-2)*180°}{n} = 360° \qquad (8)$$

The solution is:

$$x = 2 + \frac{4}{n-2} \qquad (9)$$

Because x is positive integer, then:

$$
\begin{aligned}
n &= 3, x_3 = 6; \\
n &= 4, x_4 = 4; \\
n &= 6, x_6 = 3;
\end{aligned} \qquad (10)
$$

From that we can see the polygon can be regular triangle, square and the hexagon like the following Figure 2-1 shows.

（ a ） 　　（ b ） 　　（ c ）

Figure 2-1. three different kinds of repetitive units

Locate the nodes according to the above way, then we can get the covered area without blind points.

$$A_C = A_S - A_L = \pi R^2 - \frac{nR^2}{2} \sin \frac{2\pi}{n} \qquad (11)$$

Where:

A_C means the size of overlapping area.

A_L means the size of the circle with the effective communication radius.

A_S means the size of inscribed polygon. n is the number of the sides on the polygon.

When $n \to \infty$, $A_C \to 0$. So in order to make sure, there is no blind node. When $n = 6$, the overlapping area is minimum:

$$A_{C\min} = (\pi - \frac{3\sqrt{2}}{2})R^2 \qquad (12)$$

Using the hexagon to locate the nodes is the best layout(Zhang Chaohui,2010).

Using the Definition 1 in the model one, we can calculate the coverage rate η_S:

$$\eta_S = \frac{A_S}{A_L} = \frac{3\sqrt{3}R^2/2}{\pi R^2} = 82.7\% \qquad (13)$$

Then the calculation function is:

$$g_1 = \frac{1}{\eta_S} - 1 = 0.2092 \qquad (14)$$

2.3　　The layout of nodes in *k*-coverage network

During the research on k-coverage network, the two puzzles challenge us. One is that if the geometric shapes in k-coverage network is the same to the *1*-coverage network. Another one is that how to calculate the distance between the two nodes.

We regard the second question as the breakthrough point to compute the distance between two nodes in k-coverage network. we have proven that there are three possibilities: regular triangle, square and the hexagon. And the node is located on the vertex.

We can see the vertex is the most difficult point to be covered. So if the vertex is covered by k times, the area is covered by k times.

We select one of the vertex *M*, to solve (x_M, y_M) using the equation of the

straight lines: L_1、 L_2. In order to make sure M is covered k times, which is that there are k nodes lie in the effective communication circle of M:

$$\sqrt{(x_M - x_N)^2 + (y_M - y_N)^2} \leq R^2 \qquad (15)$$

(x_N, y_N) is one of nodes N surrounded the node M.

Define a distance function:

$$G = R^2 - \sqrt{(x_M - x_N)^2 + (y_M - y_N)^2} \qquad (16)$$

Only if G>0,the target area is k-coverage.

Assume that: N is the point of intersection between L_1', L_2',

Then we can get the L_1', L_2' through translating L_1、 L_2 ,passing k_1 and

k_2 units.

Then $N(x_N, y_N)$ can be expressed by $M(x_M, y_M)$.

$$x_N = x_M(k_1);$$
$$y_N = y_M(k_2); \qquad (17)$$

Then the distance G can be expressed by k_1 k_2.

$$G(k_1, k_2) = R^2 - \sqrt{(x_M - x_N)^2 + (y_M - y_N)^2} \qquad (18)$$

We must find out the k_1 k_2, then we can get the distance between M and N.

（1） **Solution**

Taking the *2*-coverage for example, introduce above model in detail.

We have proved there are three shapes; regular triangle, square and the hexagon showed in the following Figure 2-2:

（ a ） （ c ） （ c ）
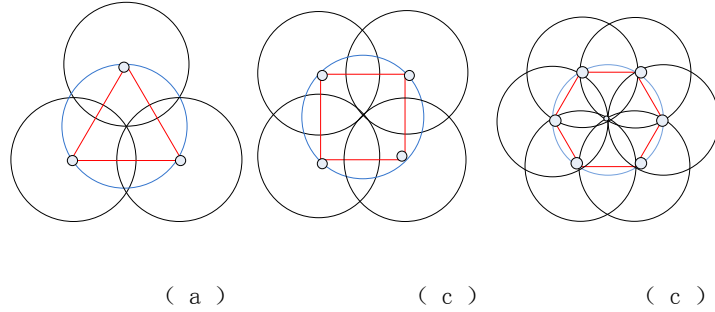
Figure 2-2. the different shapes

From above Figure 2-2(a), we can see: when it is regular triangle, it has the least nodes

So we take the *2*-coverage using regular triangle for example.

L_1、L_2 showed in the Figure 2-3:

Figure 2-3. 2-coverage using regular triangle

The equation of L_1、L_2 :

$$y_2 = \sqrt{3}x - \frac{\sqrt{3}}{2}m \qquad (19)$$

$$y_2 = -\sqrt{3}x - \frac{\sqrt{3}}{2}m \qquad (20)$$

We select one of the vertexes $M(0, \sqrt{3}m/2)$.

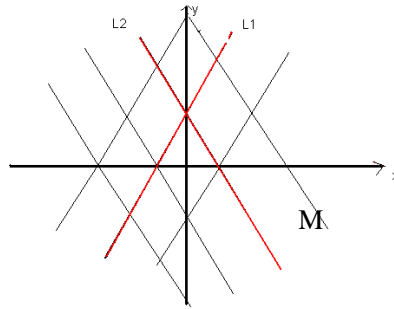Then we can get the L_1', L_2' through translating L_1、L_2 passing k_1 and

k_2 units

$$y_1 = \sqrt{3}(x + k_1 m) - m\sqrt{3}/2 \tag{21}$$

$$y_2 = -\sqrt{3}(x + k_2 m) - m\sqrt{3}/2 \tag{22}$$

Then coordinates of N:

$$x_N = -(k_1 + k_2)m/2 \tag{23}$$

$$y_N = (k_1 - k_2 - 1)\sqrt{3}m/2 \tag{24}$$

Then the distances function:

$$\begin{aligned}
G &= R_2 - (x_M - x_N)^2 + (y_M - y_N)^2 \\
&= R_2 - m_2/4*(k_1 + k_2)^2 + 3/4m^2*(k1 - k2)^2
\end{aligned} \tag{25}$$

For k–coverage, we can express k using the multiply of k_1 and k_2

It means the number of points of intersection.

$$k = ab \tag{26}$$

Where:

$$k_1 = a, k_2 = b. \tag{27}$$

From that, we can get the numerical value. Then we can get the function:

$G(k_1, k_2)$

$$\begin{aligned}
G &= R^2 - (x_M - x_N)^2 + (y_M - y_N)^2 \\
&= R^2 - m^2/4*(k_1 + k_2)^2 + 3/4m^2*(k_1 - k_2)^2 \\
&= R^2 - m^2/4*(k_1^2 - k_1 k_2 + k_2^2)
\end{aligned} \tag{28}$$

To simply the distance $G(k_1,k_2)$ into the $D(k_1,k_2)$:

$$D = k_1^2 - k_1 k_2 + k_2^2 \tag{29}$$

Because k_1 and k_2 are all integer, we can get the value of them through gradual enlargement from (0,0), for example: （0,1）,（0,-1）,（-1,0）…… After getting the value of D, put the value in order of from small to bigger. The value order is the order to pick up k_1,k_2.

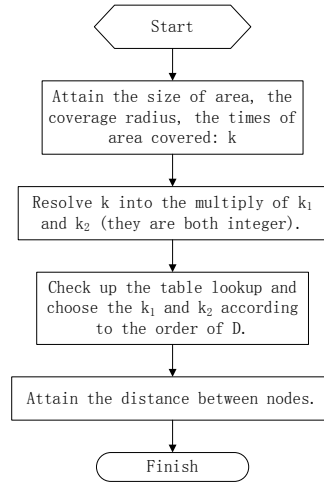We can get the value of the function D showed in Table 2-1:

Table 2-1: the value of the function D when k_1,k_2 change from -10 to 10

k1 \ k2	-10	-9	-8	-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7	8	9	10
-10	100	91	84	79	76	75	76	79	84	91	100	111	124	139	156	175	196	219	244	271	300
-9	91	81	73	67	63	61	61	63	67	73	81	91	103	117	133	151	171	193	217	243	271
-8	84	73	64	57	52	49	48	49	52	57	64	73	84	97	112	129	148	169	192	217	244
-7	79	67	57	49	43	39	37	37	39	43	49	57	67	79	93	109	127	147	169	193	219
-6	76	63	52	43	36	31	28	27	28	31	36	43	52	63	76	91	108	127	148	171	196
-5	75	61	49	39	31	25	21	19	19	21	25	31	39	49	61	75	91	109	129	151	175
-4	76	61	48	37	28	21	16	13	12	13	16	21	28	37	48	61	76	93	112	133	156
-3	79	63	49	37	27	19	13	9	7	7	9	13	19	27	37	49	63	79	97	117	139
-2	84	67	52	39	28	19	12	7	4	3	4	7	12	19	28	39	52	67	84	103	124
-1	91	73	57	43	31	21	13	7	3	1	1	3	7	13	21	31	43	57	73	91	111
0	100	81	64	49	36	25	16	9	4	1	0	1	4	9	16	25	36	49	64	81	100
1	111	91	73	57	43	31	21	13	7	3	1	1	3	7	13	21	31	43	57	73	91
2	124	103	84	67	52	39	28	19	12	7	4	3	4	7	12	19	28	39	52	67	84
3	139	117	97	79	63	49	37	27	19	13	9	7	7	9	13	19	27	37	49	63	79
4	156	133	112	93	76	61	48	37	28	21	16	13	12	13	16	21	28	37	48	61	76
5	175	151	129	109	91	75	61	49	39	31	25	21	19	19	21	25	31	39	49	61	75
6	196	171	148	127	108	91	76	63	52	43	36	31	28	27	28	31	36	43	52	63	76
7	219	193	169	147	127	109	93	79	67	57	49	43	39	37	37	39	43	49	57	67	79
8	244	217	192	169	148	129	112	97	84	73	64	57	52	49	48	49	52	57	64	73	84
9	271	243	217	193	171	151	133	117	103	91	81	73	67	63	61	61	63	67	73	81	91
10	300	271	244	219	196	175	156	139	124	111	100	91	84	79	76	75	76	79	84	91	100

The constructor can refer to this application manual table lookup to attain the numerical value of k_1 and k_2, then they can get the best layout. Using the same solutions, we can get the results of square and the hexagon.

（2）Algorithm

The flow chart in k-coverage network,See Figure 2-4:

```
                    ⟨   Start   ⟩
                         │
         ┌───────────────────────────────┐
         │ Attain the size of area, the   │
         │ coverage radius, the times of  │
         │    area covered: k             │
         └───────────────────────────────┘
                         │
         ┌───────────────────────────────┐
         │ Resolve k into the multiply of k₁ │
         │ and k₂ (they are both integer).   │
         └───────────────────────────────┘
                         │
         ┌───────────────────────────────┐
         │ Check up the table lookup and  │
         │ choose the k₁ and k₂ according │
         │    to the order of D.          │
         └───────────────────────────────┘
                         │
         ┌───────────────────────────────┐
         │ Attain the distance between nodes. │
         └───────────────────────────────┘
                         │
                    (   Finish   )
```

Figure 2-4. flow chart in k-coverage network

（3）**Cost of network**

Definition 1 Minimum repeat unit: the minimum unit while attaching the nodes, namely, triangular, square and hexagonal.

Definition 2 Node densities: the ratio of minimum repeat units and the number of the nodes in it. Symbol by n.

Because the cost of network only related to the numbers of nodes, the cost of a network can be expressed by S/n, where S represent the area of minimum repeat unit.

Table 2-2: the value of n

scheme	n
triangular	1/2
square	3/4
hexagonal	2

From the three discussing above Table 2-3 , the evaluation function is

$$g = q * k * \frac{S}{n}$$

(30)

2.4   Communication Coordination

When the effective radius of repeater is larger than or equals to the area radius (some repeater radius can reach more than 100 meters), only a repeater can cover the whole area. In order to mitigate interference, besides the geographical separation, the "continuous tone-coded squelch system" (CTCSS), sometimes nicknamed "private

line" (PL), Since a repeater has a special frequency pair and a particular PL, and there are different demand of network serves (different frequencies used for different purposes) and the number of user, the spectrum range need to be divided into multiple channels.

So our task is optimizing frequency channel coordinate and the use of PL according to the users' quantities and purposes on the basis of complete coverage.

3 Results and analysis

We already know that the layout scheme of triangular, square and hexagon by model two. The graphics respectively shown as figure…

Hypothesis environmental conditions and sensor performance is the same. Then analyze their advantages and disadvantages below Figure3-1.

| a) triangular | b) square | c)hexagon |
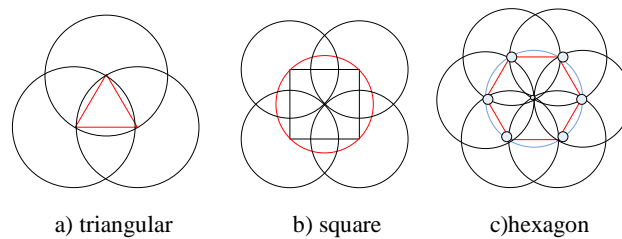
Figure3-1: Layout scheme

Use Monte Carlo Simulation to calculate the three area ratio of double covered area with the whole area. The result shows in table3-1and table3-2.

Table 4-1: The area ratio in square

Simulation times	The number of nodes	The area ratio
50000	268	0.00536
80000	441	0.005513
100000	554	0.00554
500000	2709	0.005418
900000	4989	0.005543

Table 3-2: The area ratio in hexagon

Simulation times	The number of nodes	The area ratio
50000	14678	0.29356
80000	23475	0.293438
100000	29393	0.29393

500000	146564	0.293128
900000	263982	0.293313

We have known that the double covered network in triangular scheme is actually covered triple by model two. So the ratio of double covered area with the whole area is 0 that is the positioning accuracy is 1.

As for square and hexagon, we take the area ratio by most times of simulation to calculate positioning accuracy and function values. The results show in table3-3:

Table 3-3: The positioning accuracy and function values

scheme	triangular	square	hexagon
Area ratio	0	0.005543	0.293313
positioning accuracy	1	1	0.71
function value	$0.866 \ q*R^2$	$1.333 \ q*R^2$	$0.922 \ q*R^2$

The higher value of function is, the better scheme of layout is. In order to select better layout scheme, we give a weight to each scheme based on the environment as table3-4 shows.

Table 3-4: The weights based on environment

environment	Good	Medium	bad
triangular	0--0.3	0.3-0.7	0.7--1
square	0.7--1	0.3-0.7	0--0.3
hexagon	0--0.3	0.7--1	0.3-0.7

In conclusion, the layout scheme is as follows.

(1) If the environment condition is good, choose the square scheme.
(2) If the environment condition is medium, choose the hexagon scheme.
(3) If the environment condition is bad, choose the triangular scheme.

From table 3-4, The positioning accuracy of triangular and square scheme can reach to 1, but is too low in hexagon scheme. So we take the method of reducing the distance of nodes to increase the positioning accuracy.

Use Monte Carlo Simulation; we get the relation of nodes distance and the positioning accuracy. It's show in table3-5:

Table 3-5: The relation of nodes distance and the positioning accuracy

Nodes distance	positioning accuracy

$m = R$	71%
$m = 0.9R$	81%
$m = 0.8R$	89.30%
$m = 0.7R$	95.90%
$m = 0.6R$	99.70%
$m = 0.5R$	100%

So the procedure of layout is to choose the scheme based on the environment. And if the hexagon scheme is chosen, choose the nodes distance based on positioning accuracy.

4　Discussions

(1) To once complete coverage network, choose the hexagon scheme, and the node lies in the center, the length of sides is the effective radius.

(2) To more than once complete coverage, compared the triangular, square and hexagon scheme, and get a general scheme of multiple complete coverage by calculate the length of sides on the polygon.

(3) To locating network, compared the advantages and disadvantages of triangular, square and hexagon scheme based on positioning accuracy of double complete coverage.

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